

# Extensible Program Transformation Systems

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## Mission Statement of the Stratego/XT Project

To design a language and toolset for expressing all aspects of all kinds of program transformations for all kinds of languages, concisely and efficiently.

## Transformations

- Simplification
- Optimization
  - data-flow transformation
  - inlining
  - partial evaluation
  - vectorization
  - ...
- Assimilation
- Generation
- Compilation
- Refactoring
- Documentation generation
- ...

## Languages

- Procedural
- Functional
- Object-oriented
- Logic
- Anything with structured representation

## Implementation Aspects

- Syntax definition (SDF)
- Tree representation (ATerm)
- Pretty-printing (GPP)
- Transformation strategies (Stratego)
- Transformation tool composition (XTC)

## Research Directions

- 1 Universal abstractions for program transformation

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- 2 Domain-specific language embedding
- 3 Extensible transformation systems
- 4 Meta-programming by program transformation  
(not discussed now)

Stratego: universal abstractions for program transformation

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### Core abstractions

- Matching and building terms
- Sequential composition and choice
- One-level traversal operators (specific and generic)
- Dynamic rules (context-sensitive rewriting)

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### Derived abstractions

- Match-build combinations
  - rewrite rules, term wrap, ...
- Full tree traversals
- Generic strategies

## Constant Folding

EvaPlus :

```
Plus(Int(i), Int(j)) -> Int(k)
where <add>(i,j) => k
```

EvaMul :

```
Mul(Int(i), Int(j)) -> Int(k)
where <mul>(i,j) => k
```

EvalBinOp =

```
EvalPlus <+ EvalMul <+ ...
```

const-fold =

```
all(const-fold); try(EvalBinOp)
```

## Constant propagation

```
(x := 3;  
y := x + 1;  
if foo(x) then  
  (y := 2 * x;  
   x := y - 2)  
else  
  (x := y;  
   y := 23);  
z := x + y)
```

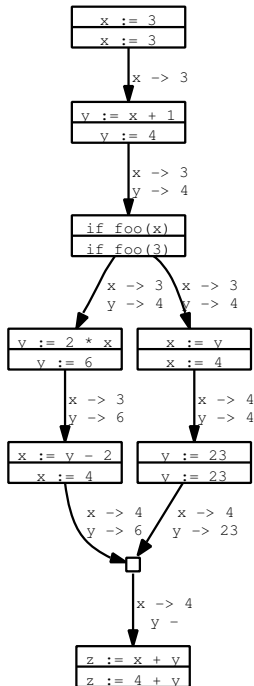
```
(x := 3;  
y := 4;  
if foo(3) then  
  (y := 6;  
   x := 4)  
else  
  (x := 4;  
   y := 23);  
z := 4 + y)
```

## Constant propagation

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(x := 3;  
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   x := 4)  
else  
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   y := 23);  
z := 4 + y)
```

using dynamic rewrite rules



## Constant Propagation Strategy

```
prop-const = PropConst <+ prop-const-assign  
  <+ prop-const-if <+ prop-const-while  
  <+ (all(prop-const); try(EvalBinOp))
```

```
prop-const-assign =  
  Assign(?x, prop-const => e)  
  ; if <is-value> e  
    then rules( PropConst : Var(x) -> e )  
    else rules( PropConst :- Var(x) ) end
```

```
prop-const-if =  
  If(prop-const, id, id)  
  ; If(id,prop-const,id) /PropConst\ If(id,id,prop-const)
```

```
prop-const-while =  
  ?While(e1, e2)  
  ; (/PropConst\* While(prop-const, prop-const))
```

## Generic Data-Flow Transformation Strategy

```
forward-prop(transform, before, after | Rs) =
let fp = prop-assign <+ prop-if <+ prop-while
      <+ transform(fp) <+ (before; all(fp); after)

prop-assign =
  Assign(id, fp)
  ; (transform(fp)
     <+ before; ?Assign(Var(x), e)
     ; undefine-dynamic-rules(|Rs,x); after)

prop-if =
  If(fp, id, id)
  ; (transform(fp) <+
     before; (If(id,fp,id) /~Rs\ If(id,id,fp)); after)
  ...
in fp end
```

## Common-Subexpression Elimination Strategy

```
cse =  
  forward-prop(fail, id, cse-assign <+ CSE | ["CSE"])
```

```
cse-assign =  
  ? Assign(Var(x), e)  
  ; where( <pure-and-not-trivial(|x)> e )  
  ; where( get-var-dependencies => xs )  
  ; rules( CSE : e -> Var(x) depends on xs )
```

```
(x := a + b;  
 y := a + b;  
 z := a + c;  
 a := 1;  
 z := (a + c) + (a + b))
```

$\Rightarrow$

```
(x := a + b;  
 y := x;  
 z := a + c;  
 a := 1;  
 z := (a + c) + (a + b))
```

## Common-Subexpression Elimination Strategy

```
cse =  
  forward-prop(fail, id, cse-assign <+ CSE | ["CSE"])  
  
cse-assign =  
  ? Assign(Var(x), e)  
  ; where( <pure-and-not-trivial(|x)> e )  
  ; where( get-var-dependencies => xs )  
  ; rules( CSE : e -> Var(x) depends on xs )
```

### Record dependencies of dynamic rules

```
(x := a + b;  
 y := a + b;  
 z := a + c;  
 a := 1;  
 z := (a + c) + (a + b))
```

$\Rightarrow$

```
(x := a + b;  
 y := x;  
 z := a + c;  
 a := 1;  
 z := (a + c) + (a + b))
```

## Domain-specific language embedding

### MetaBorg

- Embedding (domain-specific) language in host language
- Assimilation of embedded code with host code

### General architecture

- Modular syntax definition with SDF
- Assimilation with Stratego
- Combine arbitrary languages

## Concrete Syntax in Stratego

Use concrete syntax of object language in term patterns

Normalize :

```
|[ for x := e1 to e2 do e3 ]| ->  
  |[ for y := 1 to (e2 - e1) + 1 do e3 ]|  
where <not(?|[ 1 ]|)> e1  
  ; new => y  
  ; rules( ReplIV : |[ x ]| -> |[ y + e1 - 1 ]| )
```

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where <not(?|[ 1 ]|)> e1
  ; new => y
  ; rules( ReplIV : |[ x ]| -> |[ y + e1 - 1 ]| )
```

Normalize :

```
For(Var(x), e1, e2, e3) ->
  For(Var(y), Int("1"), Plus(Minus(e2, e1),Int("1")), e3)
where <not(?Int("1"))> e1; new => y
  ; rules(
    ReplIV : Var(x) -> Minus(Plus(Var(y),e1),Int("1"))
  )
```

```

public class HelloWorld {
    public static void main(String[] ps) {

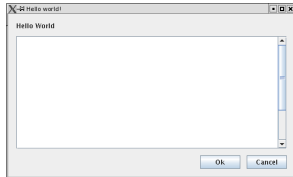
        JPanel panel = panel of border layout {
            north = label "Hello World"

            center = scrollpane of textarea {
                rows      = 20
                columns   = 40
            }

            south = panel of border layout {
                east = panel of grid layout {
                    row = {
                        button "Ok"
                        button "Cancel"
                    }
                }
            }
        }
    };
    ...
}

```

## JavaSwul



```
public class HelloWorld {
public static void main(String[] ps) {

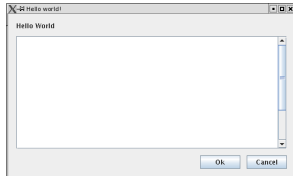
    JTextArea text = new JTextArea(20,40);

    JPanel panel
        = new JPanel(new BorderLayout(12,12));
    panel.add(BorderLayout.NORTH,
              new JLabel("Hello World"));
    panel.add(BorderLayout.CENTER,
              new JScrollPane(text));

    JPanel south
        = new JPanel(new BorderLayout(12,12));
    JPanel buttons
        = new JPanel(new GridLayout(1, 2, 12, 12));
    buttons.add(new JButton("Ok"));
    buttons.add(new JButton("Cancel"));

    south.add(BorderLayout.EAST, buttons);
    panel.add(BorderLayout.SOUTH, south);
    ...
}
```

## JavaSwul



## Extensible Transformation Systems: Use Cases

- Extend language
- Add transformation component

## Extensible Transformation Systems: Issues

- Language extension requires extension of all transformations
- Anticipated vs unanticipated extension
- Concurrent extensions
- Design mismatch

## Challenge

- A generic approach for extending transformation systems (instead of an extensible compiler for a specific language)
- What are the universal abstractions of extensibility?

## Current Extension Mechanisms in Stratego/XT

- Transformation plugins
- Syntax plugins
- Parameterized strategies
- Extending definitions
- Adaptation with aspects (experimental)

## Parameterized Strategies

### Advantages

- Recombine reusable items
- Concurrent extensions

### Disadvantages

- Anticipate variation points and lots of parameters
- No reuse of composition
- May not match needs of extension

```
super-opt =  
  forward-prop(prop-const-transform, bvr-before  
  , bvr-after; copy-prop-after; prop-const-after; cse-after  
  | ["PropConst", "CopyProp", "CSE"], [], ["RenameVar"])
```

## Extending Definitions

- Advantage: no parameter overhead, no recomposition
- Disadvantage: no concurrent extensions

## Base

```
typecheck = bottomup(try(typecheck-rules))
typecheck-rules = typecheck-if <+ typecheck-plus
typecheck-plus :
  |[ (e1 :: Int) + (e2 :: Int) ]| -> |[ (e1 + e2) :: Int ]|
```

## Extension

```
typecheck-rules = typecheck-panel
typecheck-panel :
  |[ panel of 1 { fld* } ]| ->
  |[ panel of 1 { fld* } :: JPanel ]|
where ...
```

## Adaptation with Aspects

```
module forward-prop-usage
imports forward-prop
aspects
  aspect prop-const =
    around : prop-rule(r) = proceed(["PropConst"])
    around : prop-rule = PropConst
    around : forward-prop = (proceed; try(EvalBinOp))
    before : prop-assign-ext =
      ?Assign(Var(x),e)
      ; if <is-value> e
        then rules(PropConst: Var(x) -> e)
        else rules(PropConst:- Var(x)) end
    around : prop-if = EvalIf; forward-prop <+ proceed
    around : prop-while =
      (While(forward-prop,id); EvalWhile <+ proceed)
```

## Summary

- Concise specification of transformation strategies
- Easy extension of languages
- How to integrate these?

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What are the universal abstractions of extensibility?