

# RI-MAC: A Receiver-Initiated Asynchronous Duty Cycle MAC Protocol for Dynamic Traffic Loads in Wireless Sensor Networks

Yanjun Sun, Omer Gurewitz, and David B. Johnson

SenSys 2008

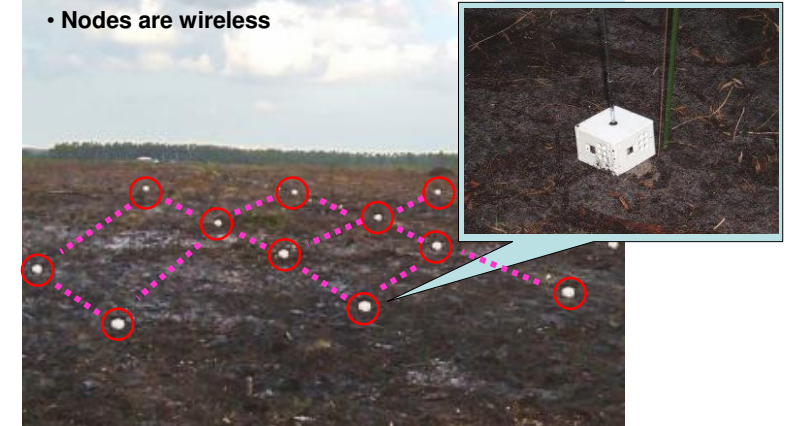


Mobile Networking Architectures



# Multihop Wireless Sensor Networks

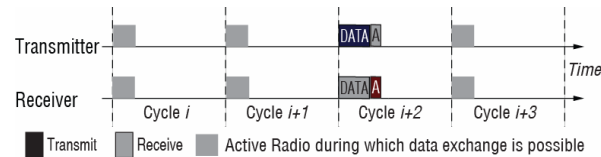
- Nodes are battery powered
- Nodes are wireless



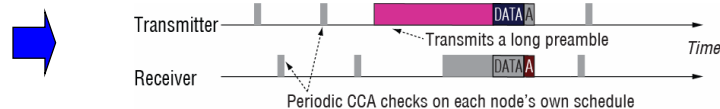
Figures from <http://cast.cse.ohio-state.edu/exscal>

# Duty Cycle MAC Protocols

- Synchronous protocols (e.g., S-MAC, T-MAC, DW-MAC)



- Asynchronous protocols (e.g., B-MAC and X-MAC)



- Hybrid protocols (e.g., WiseMAC and SCP)

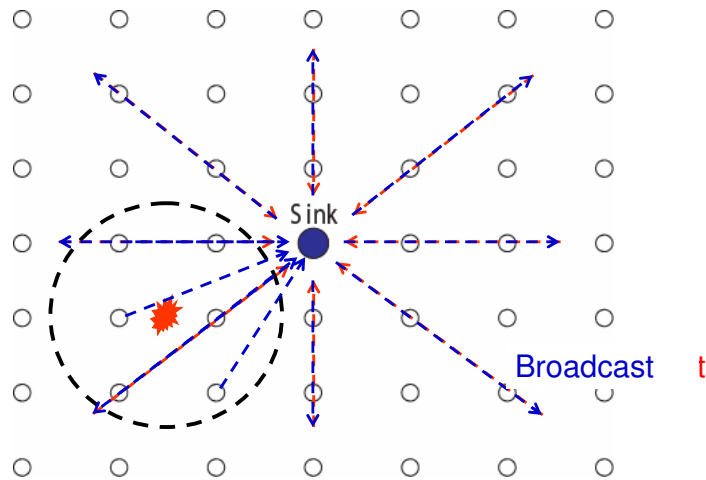


# Objectives

- Asynchronous duty cycle MAC protocol
- High power efficiency
- High packet delivery ratio
- Low packet delivery latency
- Efficient under a wide range of traffic patterns and loads

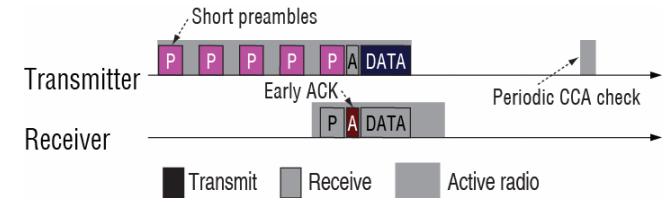


## Dynamic Traffic in Wireless Sensor Networks

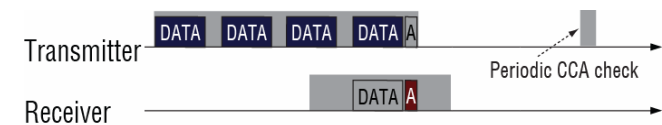


## Related Work

- X-MAC, uses **short preambles** to find rendezvous time

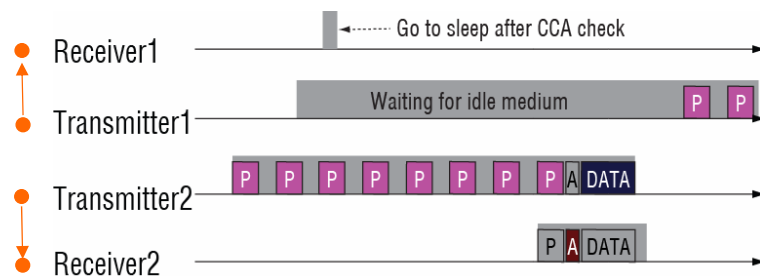


- X-MAC in TinyOS UPMA package (X-MAC-UPMA)



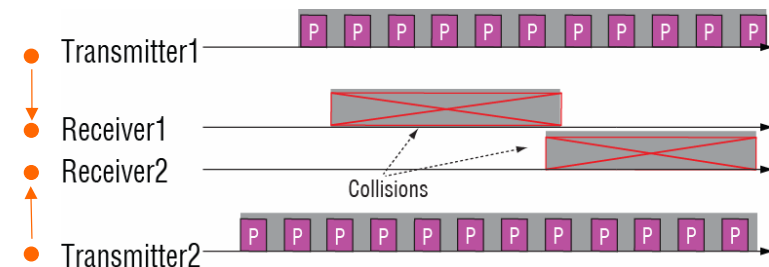
## Problems

- This preamble occupies the medium for much longer than actual data
- Blocks transmissions from neighbors



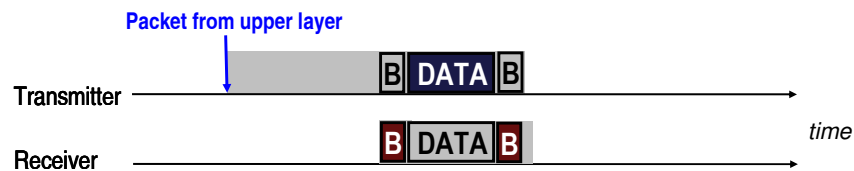
## Problems

- Preambles from hidden nodes continue colliding for a long time
- Wastes time/energy, accomplishes nothing



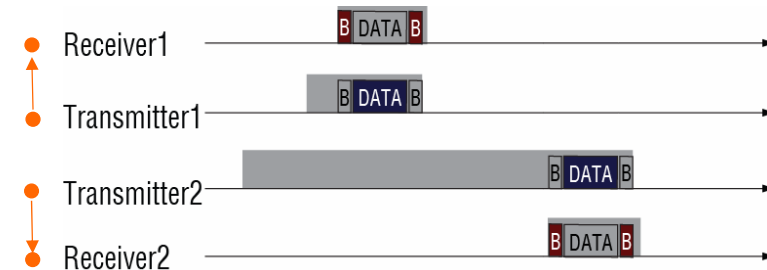
## RI-MAC: Receiver-Initiated MAC

- Transmitter wakes up and waits for the intended receiver
- Each node periodically wakes up and broadcasts a beacon after sensing idle medium
  - Sleep intervals randomly chosen from  $[0.5 \times \text{interval}, 1.5 \times \text{interval}]$
- Upon receiving a beacon from intended receiver, the transmitter starts DATA transmission
- Receiver acknowledges the DATA with another beacon



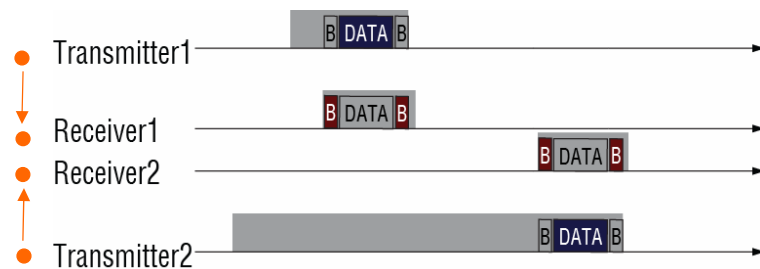
## RI-MAC Solves Earlier Problems

- Improves channel utilization
  - Reduces the time a pair of nodes occupy the medium to reach a rendezvous time



## RI-MAC Solves Earlier Problems

- Reduces collisions caused by hidden nodes

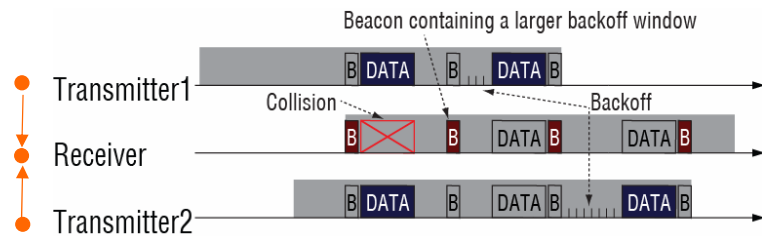


## Receiver-Controlled Backoff Window

- Beacon specifies backoff window size to be used by a transmitter before DATA transmission
  - Transmitter picks backoff delay over this range
- Receiver increases the backoff window size in beacons using binary exponential backoff

## Receiver-Controlled Backoff Window

- Reduces collision recovery cost
  - A receiver detects collisions quickly and controls retransmissions



## RI-MAC Evaluation

- ns-2* simulation-based evaluation:
  - Clique networks
  - Grid networks
  - Random networks
- TinyOS-based implementation and evaluation:
  - Clique networks
  - Networks with hidden nodes

## Simulation Parameters

- Radio parameters

Bandwidth	250 Kbps	Size of Hardware Preamble	6 B
SIFS	192 $\mu$ s	Size of ACK	5 B
Slot time	320 $\mu$ s	CCA Check Delay	128 $\mu$ s
Tx Range	250 m	Carrier Sensing Range	550 m

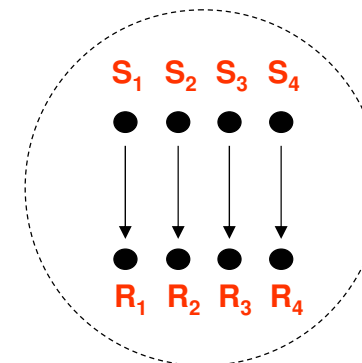
- MAC protocol parameters

	XMAC	XMAC-UPMA	RI-MAC
Backoff Window	32	32	0–255
Retry Limit	0 of 5	0 of 5	5
Special Frame	Short Preamble	—	Beacon
Special Frame Size	6 B	—	6–9 B
Dwell Time	10.5 ms	100 ms	Variable

- 1-second duty cycle interval; 28-byte DATA frames

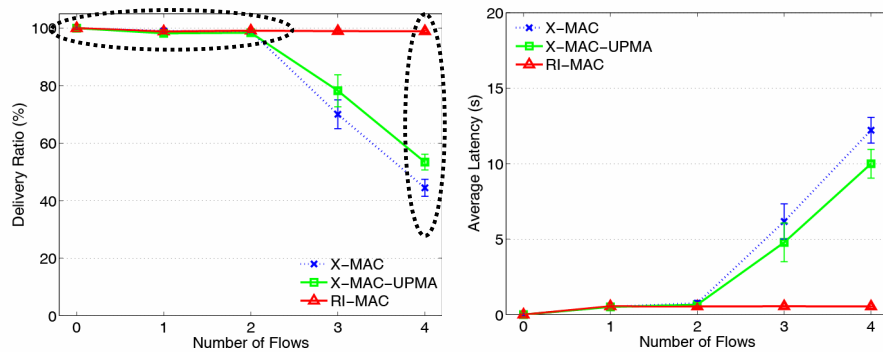
## Clique Networks

- Independent flows that do not share end nodes
- Packet generation interval is randomly chosen from [0.5, 1.5] seconds for each flow



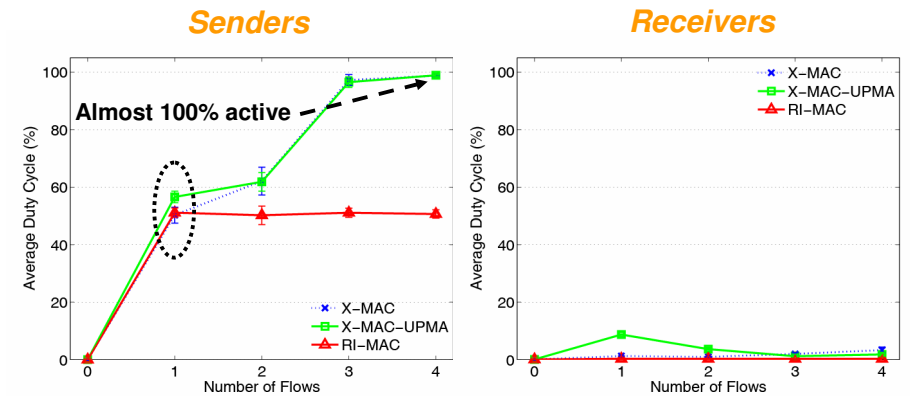
## Clique Networks: PDR and Latency

- Independent flows that do not share end nodes
- Packet generation interval is randomly chosen from [0.5, 1.5] seconds for each flow



## Clique Networks: Average Duty Cycle

Average Duty Cycle = Percentage of time a node is on



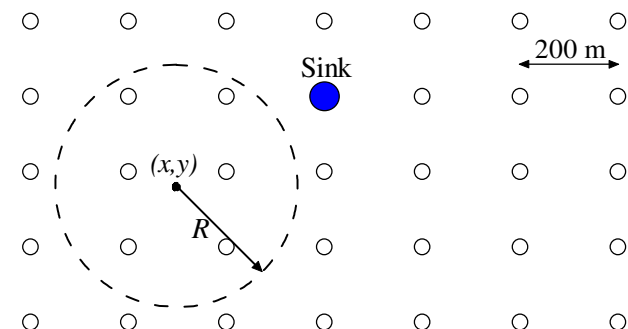
## Grid Network

- 49 nodes in  $7 \times 7$  grid
- Sink node is at center node
- 10 different simulation runs
- Random correlated-event traffic with event every 100 seconds

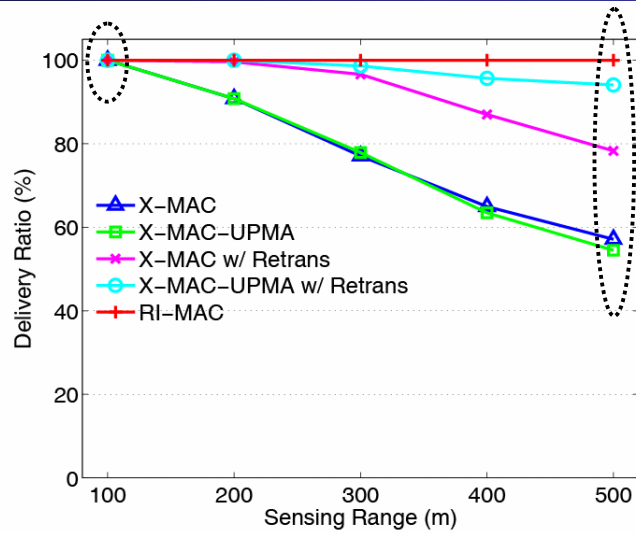
## Random Correlated-Event Traffic

Average number of packets per event vs. sensing range (R):

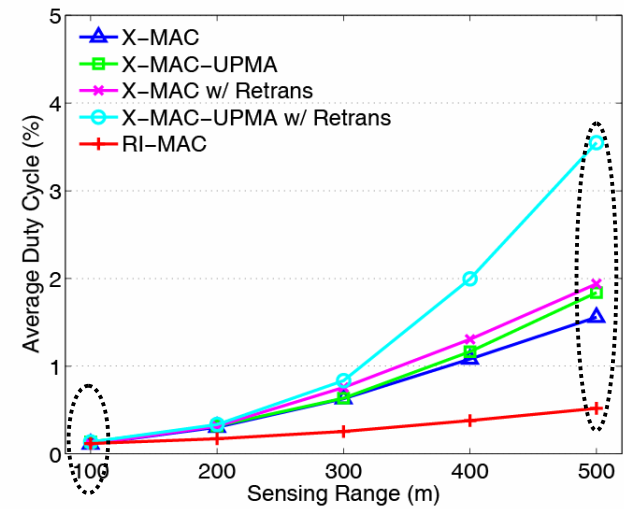
Range(m)	100	200	300	400	500
Packets	0.8	3.1	6.4	10.6	15.2



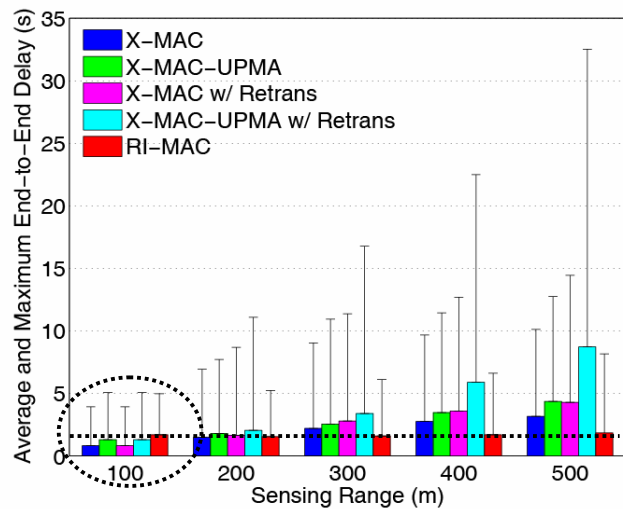
## Grid Network: Packet Delivery Ratio



## Grid Network: Average Duty Cycle



## Grid Network: End-to-End Latency

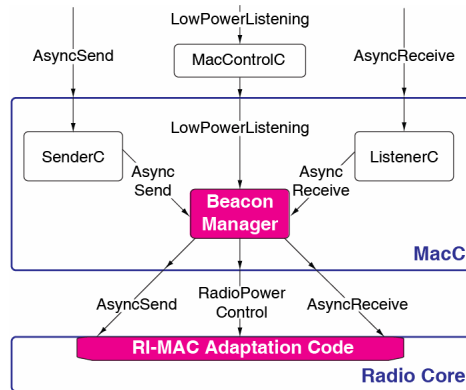


## Random Networks

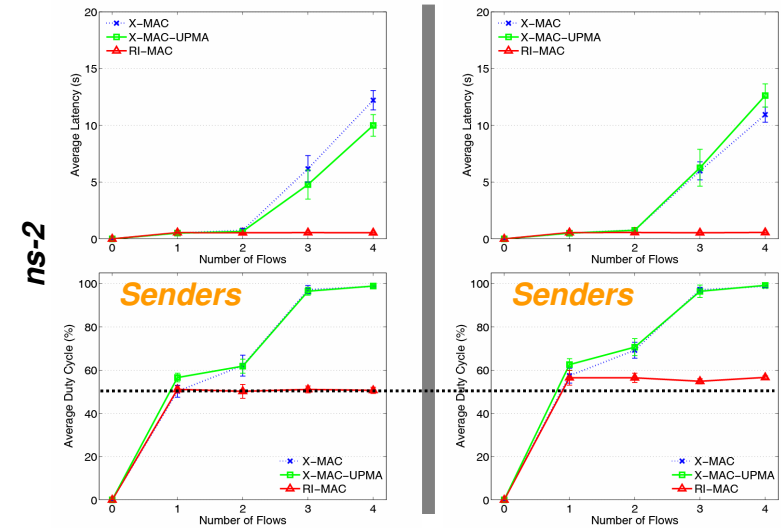
Results for random network show similar trends and are reported in the paper

## RI-MAC Implementation in TinyOS on MICAz

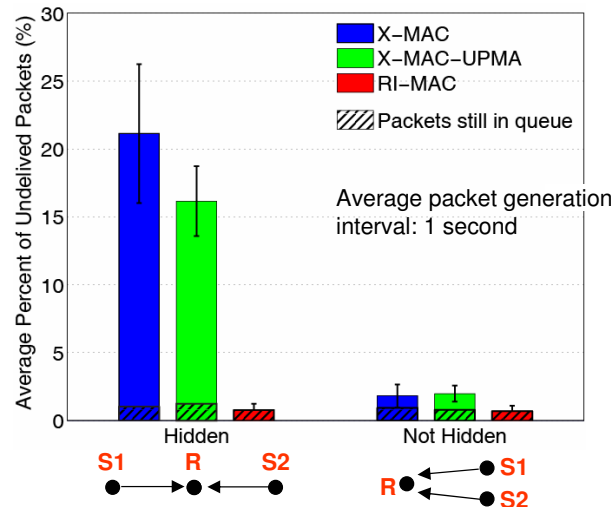
- Under the UPMA framework in TinyOS
- All beacons operations are in software



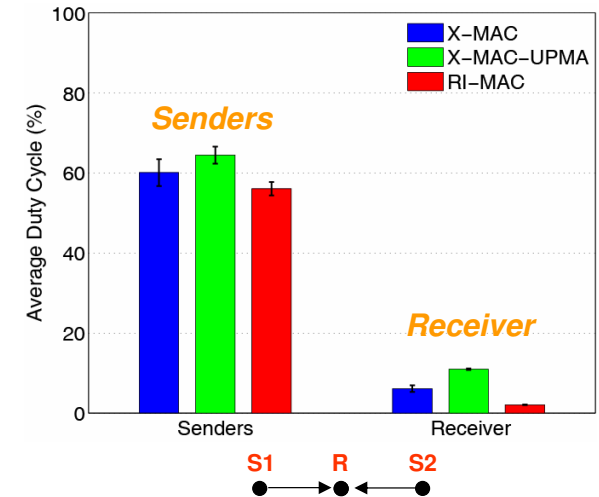
## Clique Network: Simulation vs. Experimental



## 3-Node Network with Two Hidden Nodes



## 3-Node Network with Two Hidden Nodes



## Conclusion

- ❑ RI-MAC is an asynchronous duty-cycle MAC
- ❑ Receiver-initiated transmissions and receiver controlled collision resolution
- ❑ Improved channel utilization
- ❑ Lower collision recovery cost
- ❑ Implemented in TinyOS
- ❑ Evaluated in *ns-2* and TinyOS/MICAz
- ❑ Performs much better than X-MAC and X-MAC-UPMA