Comp 411
Principles of Programming Languages
Lecture 18
Run-time Environment Representations II

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Review

• In Algol-like languages, the collection of environments that exist at any point during a computation is embedded in the machine control stack supporting (recursive) procedure calls. When the frames of the control stack are used in this way, they are called activation records.

• In each activation record, a pointer called the static link points to the environment parent of the record. Similarly, a pointer called the dynamic link points to the preceding stack frame (activation record) to which control will return when the current computation [conducted using the current activation record] completes. The static link is used for looking up non-local bindings (of free variables in the current let or lambda-abstraction) in the environment.

• The dynamic link is used to return control from the current “procedure” to its caller (whose local variables may not be accessible from the current frame).
Example I

Consider the following Scheme program to reverse a list:

```scheme
(define rev (lambda (l)
    (letrec
        [(revhelp ; :=
            (lambda (tl acc)
                (if (empty? tl) acc
                    (revhelp (rest tl) (cons (first tl) acc))))]
        (revhelp l empty))))
```

The Pidgin Algol equivalent (extended to include functional lists as built-in type):

```algol
List rev(l: List) = {
    { List revhelp(tl: List, acc: List) = {
        if empty?(tl) then acc else revhelp(rest(tl), cons(first(tl), acc)) }
        revhelp(l, empty)
    }
}
```

What happens when `(rev '(0 1))` is called?

- The top level call on `rev` allocates activation record (AR) #1 with null static and dynamic links and a slot for `l` (el) initialized to `'(0 1)`.  
- The body of `rev` (executing in AR #1) allocates AR #2 for the `letrec` with static and dynamic links pointing to preceding activation record and a slot for `revhelp` initialized to the closure for its definition.  
- The body of `revhelp` allocates AR #3 record for the recursive call on `revhelp` with static link taken from closure binding of `revhelp` (in AR #2) and dynamic link pointing to preceding activation record.
Since \( l \) is not empty, body of \texttt{revhelp} allocates AR #4 for the recursive call on \texttt{revhelp} with static link taken from closure binding of \texttt{revhelp}, dynamic link \( \ldots \), and slots for \texttt{tl} and \texttt{acc} initialized to \texttt{'(1)} and \texttt{'(0)}, respectively.

Since \( l \) is not empty, body of \texttt{revhelp} allocates AR #5 record for recursive call on \texttt{revhelp} with static link taken from closure binding of \texttt{revhelp}, dynamic link \( \ldots \), and slots for \texttt{tl} and \texttt{acc} initialized to \texttt{'() and '}(1 \ 0), respectively.

Since \( l \) is empty, body of \texttt{revhelp} in context of AR #5 returns the value \texttt{'(1 \ 0)}, popping AR #5 off the stack.

The pending evaluation in AR #4 returns the value \texttt{'(1 \ 0)}, popping AR #4.

The pending evaluation in AR #3 returns the value \texttt{'(1 \ 0)}, popping AR #3.

The pending evaluation in AR #2 returns the value \texttt{'(1 \ 0)}, popping AR #2.

The pending evaluation in AR #1 returns the value \texttt{'(1 \ 0)}, popping AR #1.

Notes:
1. The last four steps are trivial because they are returns from tail calls.
2. The dynamic link is \textit{always} set to point to the preceding AR.
3. Algol 60 was designed so that the ARs could be stack allocated (and deallocated). Function values are not “first-class”.
4. Guy Steele’s heap allocation “hack” (to be covered later) extends the Algol stack allocation runtime to support to languages supporting first-class functions/procedures.
5. In Java, inner classes enable the nesting of scopes as in Algol; the static chain is formed by embedding hidden parent instance pointers in the inner class objects. In addition all non-local variables accessed in an inner class must be final so that they can be copied into the inner class instances.
Example II

Consider the following Scheme program to reverse a list:

(define lookup (lambda (sym env)
    (letrec
        [(lookup-help
            (lambda (env)
                (cond [(empty? env) null]
                [(eq? sym (pair-var (first env))
                    (pair-val (first env))]
                [else (lookup-help (rest env) tl)]))]
        (lookup-help env))))

Let’s trace the evaluation of (lookup 'a (cons (make-pair 'a 5) null))

- The top-level call on lookup allocates AR #1 with null static link and slots for sym and env initialized to 'a and '([a 5]).
- The body of lookup (executing in AR #1) allocates AR #2 for the block with the static link pointing to AR #1 and a slot for lookup-help initialized to the closure for its definition.
- The body of the block executing in AR #2 allocates AR #3 for the call on lookup-help with the static link extracted from the closure bound to lookup-help and a slot for env initialized to '([a 5]) (the value of env in the environment determined by the static link of AR #2).
- The body of lookup-help executing in AR #3 looks at env and finds a match for sym (found in the static chain in AR #1) in the first pair, namely ['a 5] and...
Exceptions were not included in Algol 60 or most of its successors (Pascal, Algol W, C). But the Algol 60 run-time stack can easily handle the Java \texttt{try/catch} construct.

How does exception handling work? The activation record must include a \texttt{catch} table for the active \texttt{catch} (assuming one exists) listing the caught exception classes (types) and their handlers (the bodies of the \texttt{catch} clauses). (A \texttt{catch} is active if control is within the corresponding \texttt{try} block.) When an exception is thrown the executing code (interpreter or compiled code) searches back through the dynamic chain—popping exited frames off the stack—to find the first matching catch clause.