HPCToolkit : Multi-platform Tools for Profile-based Performance Analysis

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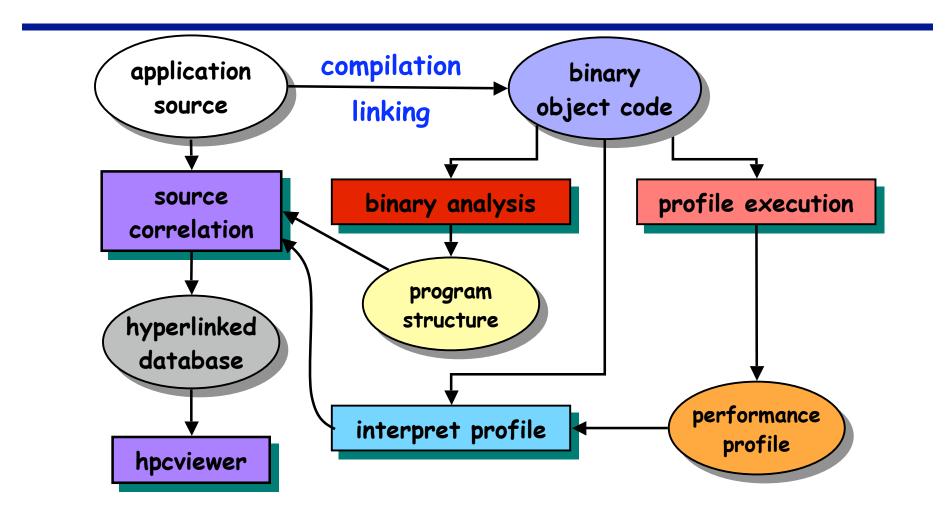


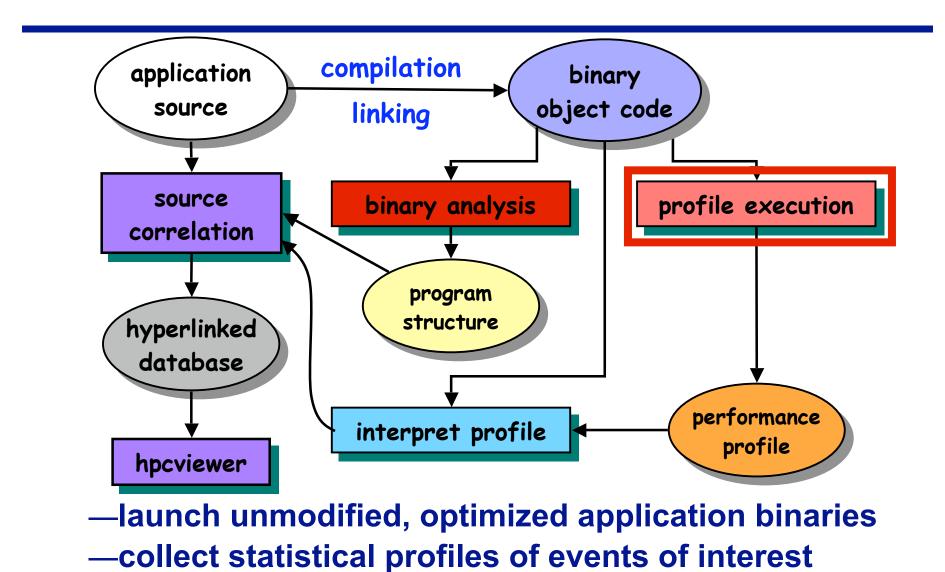
Performance Analysis and Tuning

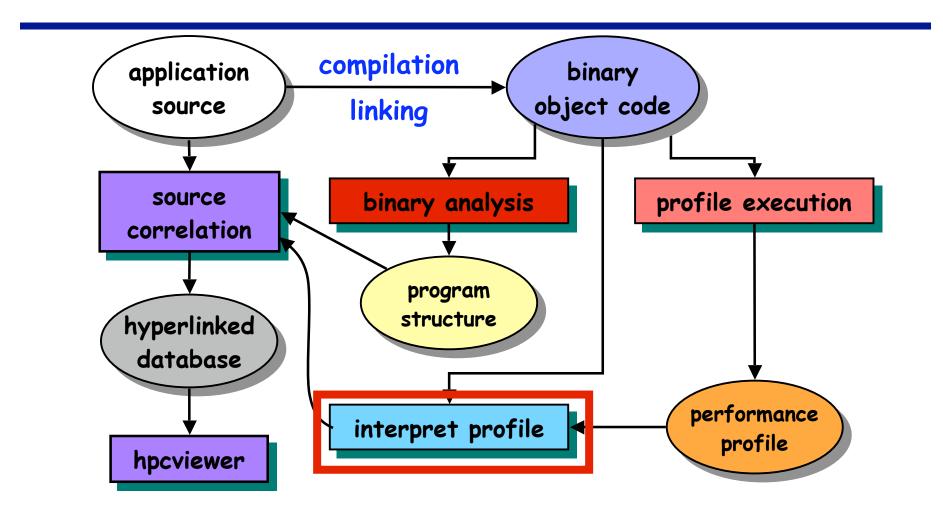
- Increasingly necessary
 - -gap between typical and peak performance is growing
- Increasingly hard
 - -complex architectures are harder to program effectively
 - complex processors
 - VLIW
 - deeply pipelined, out of order, superscalar
 - complex memory hierarchy non-blocking, multi-level caches
 - TLB
 - - multi-lingual programs
 - many source files
 - complex build process
 - external libraries in binary-only form

HPCToolkit Goals

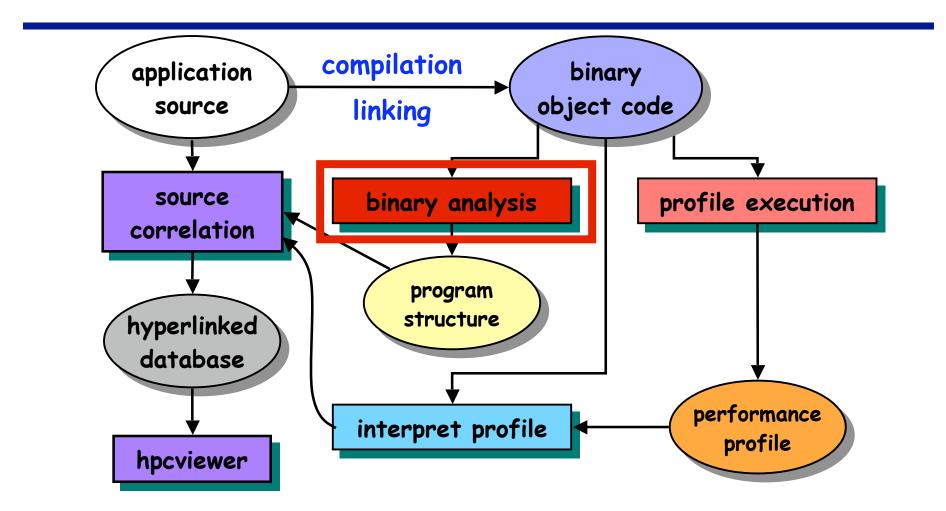
- Support large, multi-lingual applications
 - —a mix of of Fortran, C, C++
 - -external libraries
 - -thousands of procedures
 - -hundreds of thousands of lines
 - -we must avoid
 - manual instrumentation
 - significantly altering the build process
 - frequent recompilation
- Multi-platform
- Scalable data collection
- Analyze both serial and parallel codes
- Effective presentation of analysis results —intuitive enough for physicists and engineers to use —detailed enough to meet the needs of compiler writers



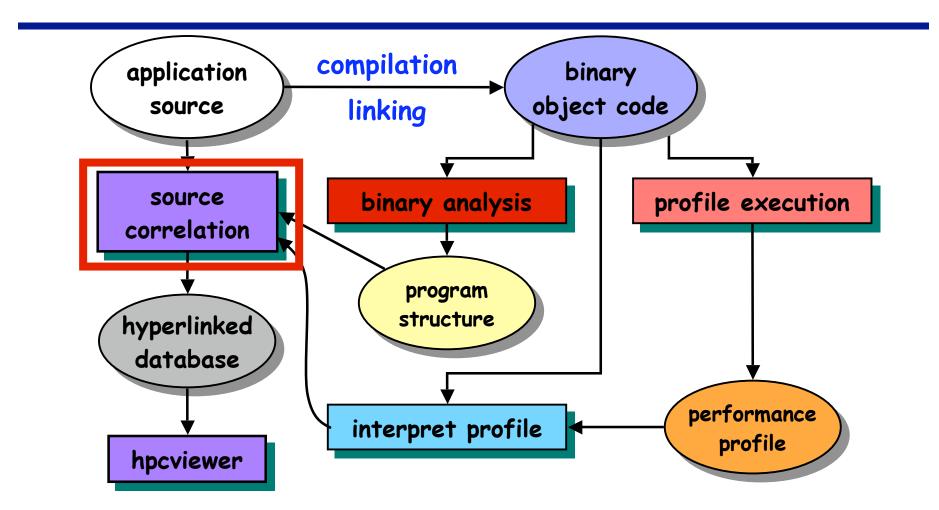




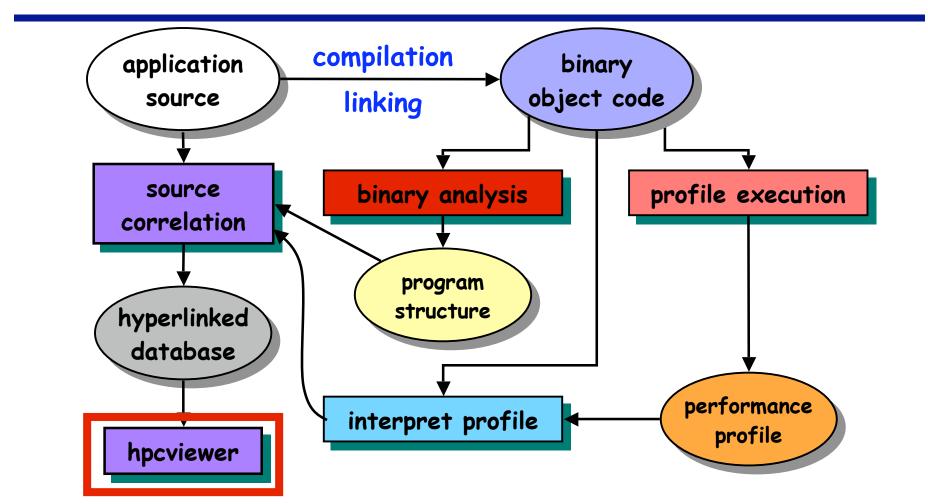
-decode instructions and combine with profile data



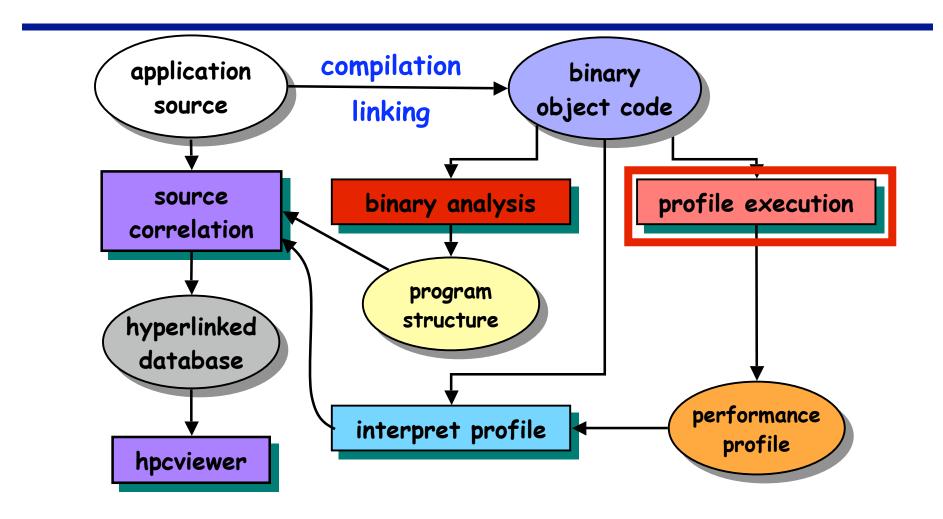
—extract loop nesting information from executables



—synthesize new metrics by combining metrics —relate metrics, structure, and program source



—support top-down analysis with interactive viewer —analyze results anytime, anywhere



Data Collection

Support analysis of unmodified, optimized binaries

- Inserting code to start, stop and read counters has many drawbacks, so don't do it!
 —nested measurements skew results
- Use hardware performance monitoring to collect statistical profiles of events of interest
- Different platforms have different capabilities

 –event-based counters: MIPS, IA64, Pentium
 –ProfileMe instruction tracing: Alpha
- Different capabilities require different approaches

Data Collection Tools

Goal: limit development to essentials only

• MIPS-IRIX:

—ssrun + prof → ptran

• Alpha-Tru64:

—uprofile + prof → ptran
—DCPI/ProfileMe → xprof

• IA64-Linux and IA32-Linux —papirun/papiprof

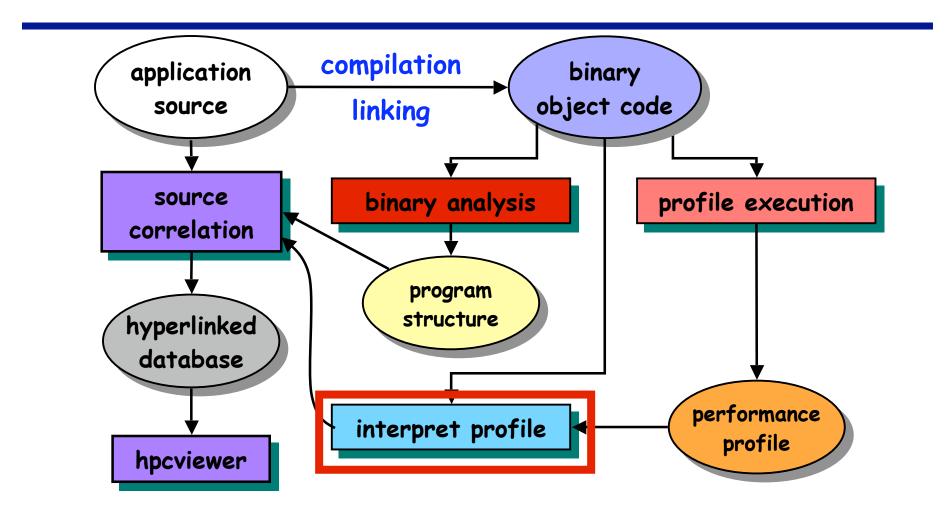
papirun/papiprof

- PAPI: Performance API

 —interface to hardware performance monitors
 —supports many platforms
- papirun: open source equivalent of SGI's 'ssrun' —sample-based profiling of an execution
 - preload monitoring library before launching application
 - inspect load map to set up sampling for all load modules
 - record PC samples for each module along with load map
 - -Linux IA64 and IA32
- papiprof: 'prof'-like tool
 - -based on Curtis Janssen's vprof
 - —uses GNU binutils to perform PC → source mapping
 - -output styles
 - XML for use with hpcview
 - plain text

DCPI and ProfileMe

- Alpha ProfileMe
 - —EV67+ records info about an instruction as it executes
 - mispredicted branches, memory access replay traps
 - more accurate attribution of events
- DCPI: (Digital) Continuous Profiling Infrastructure
 - -sample processor counters and instructions continuously during execution of *all* code
 - all programs
 - shared libraries
 - operating system
 - -support both on-line and off-line data analysis
 - to date, we use only off-line analysis

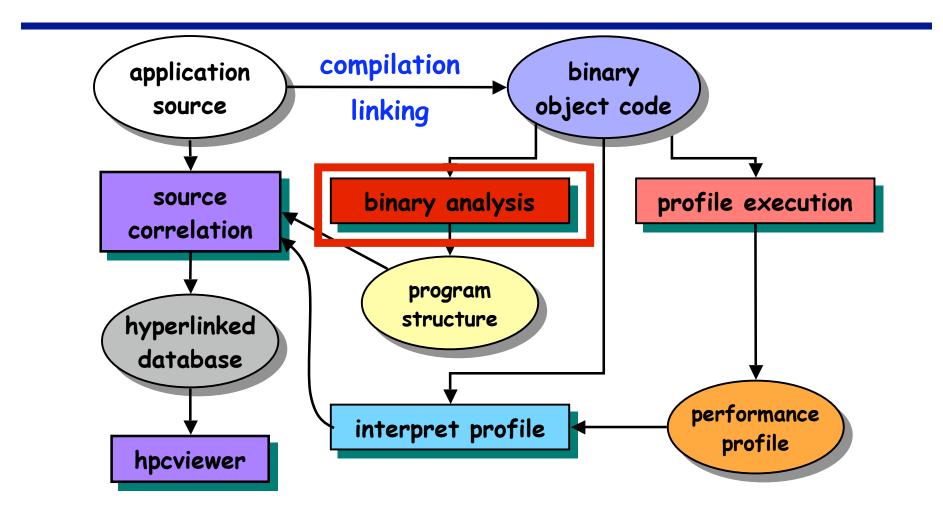


Metric Synthesis with xprof (Alpha)

Interpret DCPI samples into useful metrics

- Transform low-level data to higher-level metrics
 - —DCPI ProfileMe information associated with PC values
 - ----project ProfileMe data into useful equivalence classes
 - —decode instruction type info in application binary at each PC
 - FLOP
 - memory operation
 - integer operation
 - -fuse the two kinds of information
 - Retired instructions + instruction type = retired FLOPs retired integer operations retired memory operations

Map back to source code like papiprof

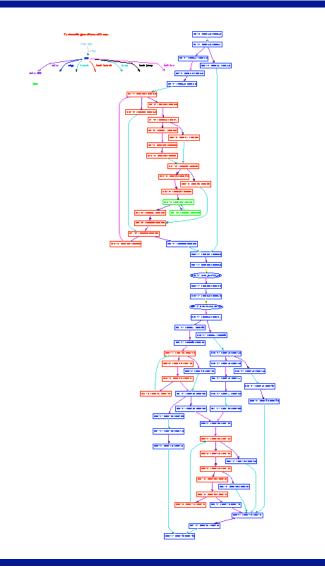


Program Structure Recovery with bloop

- Parse instructions in an executable using GNU binutils
- Analyze branches to identify basic blocks
- Construct control flow graph using branch target analysis —be careful with machine conventions and delay slots!
- Use interval analysis to identify natural loop nests
- Map machine instructions to source lines with symbol table —dependent on accurate debugging information!
- Normalize output to recover source-level view

Platforms: Alpha+Tru64, MIPS+IRIX, Linux+IA64, Linux+IA32, Solaris+SPARC

Sample Flowgraph from an Executable



Loop nesting structure —blue: outermost level —red: loop level 1 —green loop level 2

Observation optimization complicates

program structure!

Normalizing Program Structure

Constraint: each source line must appear at most once

Coalesce duplicate lines

- (1) if duplicate lines appear in different loops
 - find least common ancestor in scope tree; merge corresponding loops along the paths to each of the duplicates

purpose: re-rolls loops that have been split

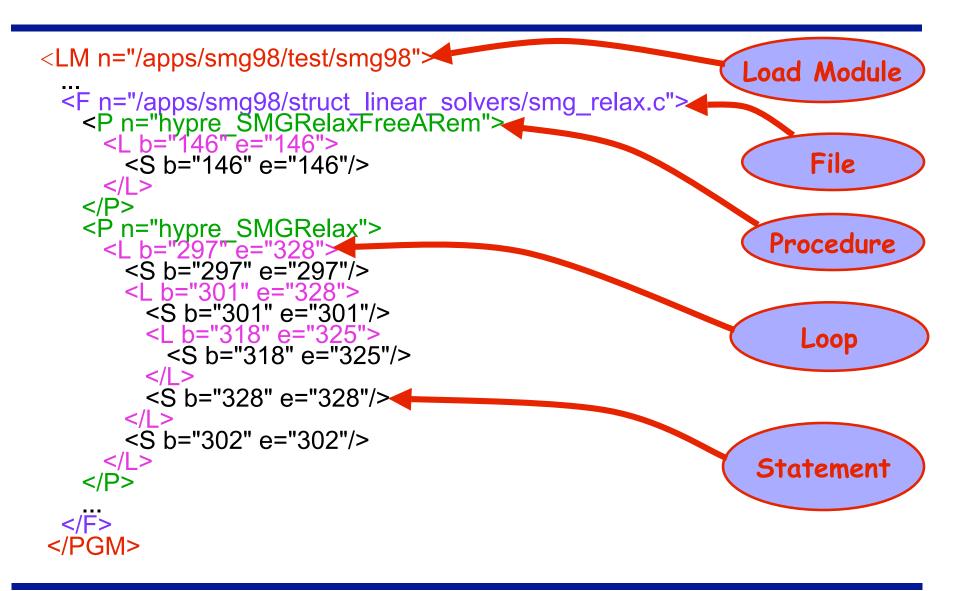
(2) if duplicate lines appear at multiple levels in a loop nest

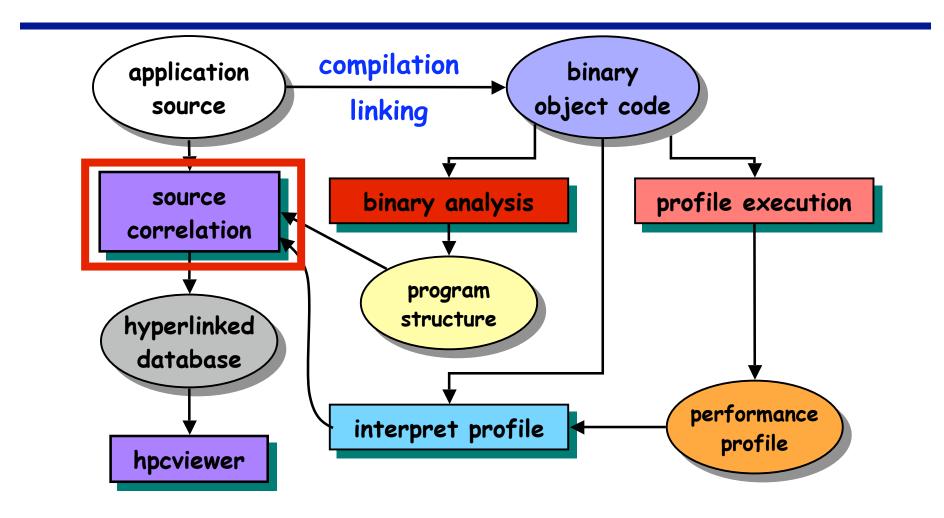
discard all but the innermost instance

purpose: handles loop-invariant code motion

apply (1) and (2) repeatedly until a fixed point is reached

Recovered Program Structure



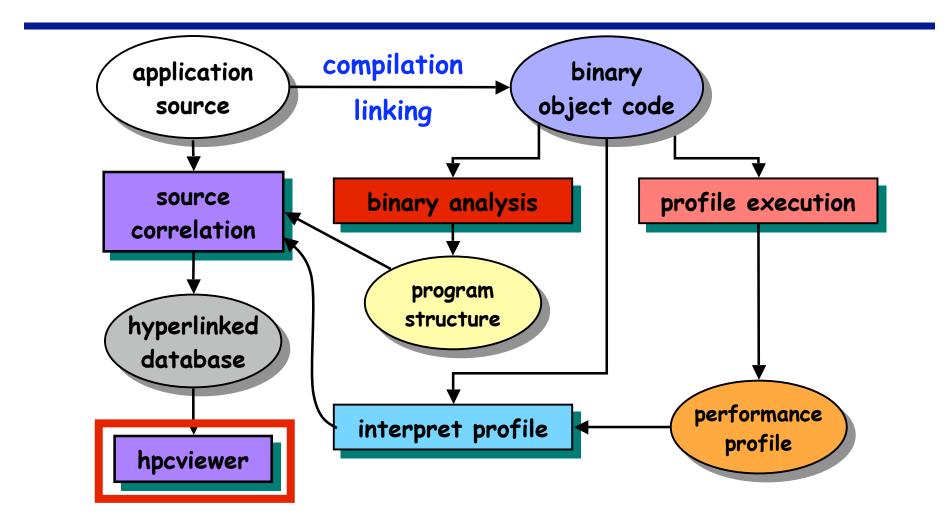


Data Correlation

Problem

—any one performance measure provides a myopic view

- some measure potential *causes* (e.g. cache misses)
- some measure effects (e.g. cycles)
- cache misses not always a problem
- —event counter attribution is inaccurate for out-of-order processors
- Approaches
 - multiple metrics for each program line
 - -computed metrics, e.g. cycles FLOPS
 - eliminate mental arithmetic
 - serve as a key for sorting
 - -hierarchical structure
 - line level attribution errors give good loop-level information

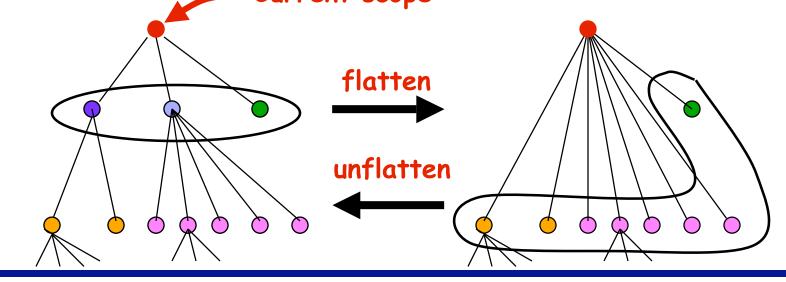


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	228	end do		HP	HPCViewer Screenshot						
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Ă	238	do $k = start(3,c), ksize$	e-end(3.c)-1								
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	242	i2 = i + 2		Annotated Source view							
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Flattening for Top Down Analysis

- Problem
 - -strict hierarchical view of a program is too rigid
 - —want to compare program components at the same level as peers
- Solution
 - —enable a scope's descendants to be flattened to compare their children as peers
 Current scope



Some Uses for HPCToolkit

Identifying unproductive work

—where is the program spending its time not performing FLOPS

- Memory hierarchy issues

 bandwidth utilization: misses x line size/cycles
 exposed latency: ideal vs. measured
- Cross architecture or compiler comparisons —what program features cause performance differences?
- Gap between peak and observed performance —loop balance vs. machine balance?
- Evaluating load balance in a parallelized code —how do profiles for different processes compare

Assessment of HPCToolkit Functionality

- Top down analysis focuses attention where it belongs —sorted views put the important things first
- Integrated browsing interface facilitates exploration —rich network of connections makes navigation simple
- Hierarchical, loop-level reporting facilitates analysis —more sensible view when statement-level data is imprecise
- Binary analysis handles multi-lingual applications and libraries —succeeds where language and compiler based tools can't
- Sample-based profiling, aggregation and derived metrics —reduce manual effort in analysis and tuning cycle
- Multiple metrics provide a better picture of performance
- Multi-platform data collection
- Platform independent analysis tool

What's Next?

Research

-collect and present dynamic content

- what path gets us to expensive computations?
- accurate call-graph profiling of unmodified executables
- analysis and presentation of dynamic content
- -communication in parallel programs
- -performance diagnosis: why rather than what

Development

- -harden toolchain
- -new platforms: Opteron and PowerPC
- -data collection with oprofile on Linux